

# SPARQL Extensibility

- Arbitrary functions in FILTERs
  - Identified by URI
  - Can extend operators as well
- New semantics for Basic Graph Patterns
  - BGPs *extract* mappings from data sets
    - The algebra is independent of the extraction
      - We hope!
  - One document/graph has many semantics
    - Simple, RDF, RDFS, OWL....
  - Queries (should be) sensitive to the semantics
    - See prior unit for some examples (RDF, etc.)

# Trickiness

- Controlling answers
  - Too many:
    - Simple entailment can yield infinite answers
  - Too few:
    - Finding all proofs of an answer difficult
- New sorts of issue
  - E.g., inconsistent data or equality
- Performance
  - Bare consistency of OWL DL is NEXPTIME
  - Query languages very expressive!
  - Performance model unclear to users

# Standardizing these things is hard

- Not a lot of experience
  - Conjunctive query for DLs is (fairly) new
    - Theoretically and implementationally
  - Concept language expressive
    - Can express many common queries
  - Lots of decisions
- Database experience not always reliable
- LP experience not always reliable

# Inconsistency

- Some logics can express **inconsistent** data
  - RDF (with certain datatypes), RDFS, and OWL
  - Inconsistencies entail **everything**
    - So, every mapping is a “**correct**” answer!
  - Inconsistencies often signal error
    - But may indicate mere disagreement!
- What should a query engine return?
  - Nothing
  - No answers, but explanations
  - Implementation dependent “best” answers
  - Answers from a weaker logic

# “Strange” queries

- In RDF(S)
  - Thin distinction between schema and data
    - Schema language very inexpressive
  - So easy to treat the schema and data uniformly
- In OWL-DL
  - Strong distinction between schema and data
    - TBox vs. Abox
    - Concept language very expressive
  - OWL Full tries to do the RDF thing
    - High cost: Undecidability, no implementations, hard to understand semantics

# Types of Query Variables

- 2 key axis of a variable with 4 combinations

- A) Distinguished
- B) "Semi-"distinguished
- C) "Projected away"
- D) Non-distinguished

**Binds to  
names  
only**

**In head of query**

A. Yes/Yes	A. Yes/No
B. No/Yes	A. No/No

- In a databases, only A and C are possible
  - C and D collapse (no non-named entities)
- In DLs, A and D are standard
  - D make query answering harder!
- In SPARQL/RDF, all variables are B

# Considerations

- A and C are easiest
  - Bind only to names
  - Reduces SPARQL over RDF to relational (mostly)
  - Tables don't contain variables
  - Lean vs. non-lean source graphs don't matter
  - Miss answers!
- Non-distinguished
  - Basic answering is much more difficult (esp. in DL)
  - But tables and algebra remain the same
  - Miss some and co-reference of answers

# Semi-distinguishedness

- Semi-distinguished

- Basic answering is much more difficult
  - Less so in RDF
  - Not clearly defined yet
    - A binding is supposed to be true in all models
    - How exactly to identify generated individuals across models?
- New issues for the algebra and final results
  - Even in RDF!



# (Non-)distinguished Example

Data:

```

: bob : loves _ : x .
_ : x : loves : bob .
_ : y : loves : sally
: sally : loves : sheevah .

```

In (pseudo-)OWL:

```

: bob a [onProperty : loves
        someValuesFrom
        [onProperty : loves;
         hasValue : bob].
: sally a [onProperty inv(: loves);
         someValuesFrom Thing]

```

SELECT ?lover	
WHERE { ?lover :loves ?beloved }	<b>?lover</b>
	:sally
WHERE { ?lover :loves _ :beloved }	<b>?lover</b>
	:bob
	:sally

# Semi-distinguished Example

- Data (co-reference in data (could) show up in results):

```
:bob :loves _:x.  
_:x :loves :bob.  
_:y :loves :sally  
:sally :loves :sheevah.
```

**SELECT ?lover ?beloved**

**WHERE { ?lover :loves ?beloved }**

<b>?lover</b>	<b>?beloved</b>
:bob	_:G1
_:G1	:bob
_:G2	:sally
:sally	:sheevah

# Oedipus Example

<pre>SELECT ?x FROM &lt;http://www.mindswap.org/ontologies/oedipus&gt;</pre>			
<pre>WHERE { ?x ns:hasChild   [ rdf:type ns:Patricide;     ns:hasChild ?y].   ?y rdf:type ns:NotPatricide.}</pre>	<p><b>NO RESULTS</b></p>		
<pre>WHERE { ?x ns:hasChild   [ rdf:type ns:Patricide;     ns:hasChild   [ rdf:type ns:NotPatricide]]}</pre>	<table border="1"> <tr> <td style="text-align: center;"><b>x</b></td> </tr> <tr> <td style="text-align: center;"><b>:IOKASTE</b></td> </tr> </table>	<b>x</b>	<b>:IOKASTE</b>
<b>x</b>			
<b>:IOKASTE</b>			

(This relies on a set of OWL DL axioms in the background.)  
 (Example from the Description Logic Handbook, chapter 2.)

# The data

```
:Patricide    a owl:Class .
:NotPatricide a owl:Class;
              owl:complementOf :Patricide .
:hasChild    a owl:ObjectProperty .

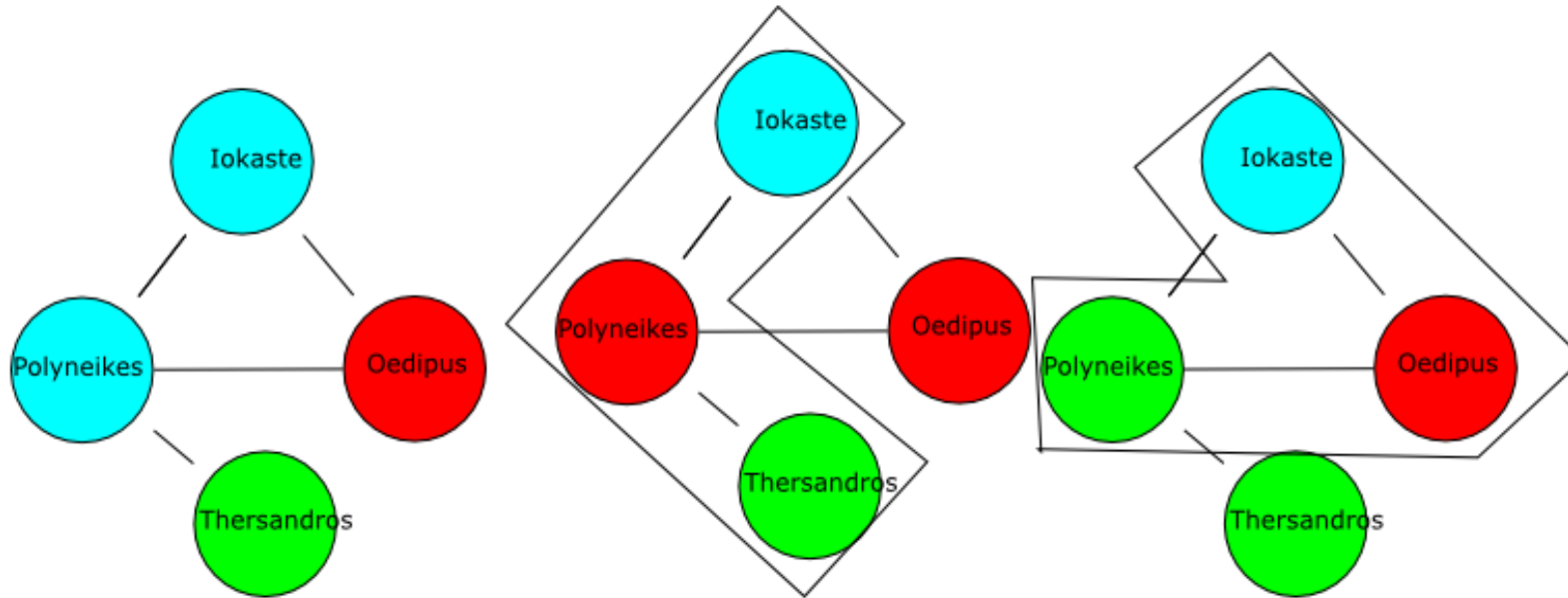
:IOKASTE
  ns:hasChild :OEDIPUS;
  ns:hasChild :POLYNEIKES .

:OEDIPUS    a ns:Patricide
  ns:hasChild :POLYNEIKES .

:POLYNEIKES
  ns:hasChild :THERSANDROS .

:THERSANDROS a ns:NotPatricide .
```

# Reasoning with Oedipus



## Reasoning by cases

- Iokaste hasChild Polyneikes, Oedipus
- Oedipus is a Parricide and Thersandros is not
- Polyneikes is either a Parricide or not a Parricide

# Thoughts and Implications

- If  $?y$  in the Oedipus query were semi-distinguished:
  - What should its value be?
  - Polyneikes or Thersandros?
  - $_:x$  instance of ( $\{Polyneikes\}$  or  $\{Thersandros\}$ )?
- In the simple/RDF case:
  - Entailment can only introduce redundant BNodes
  - The data could be non-lean
- Suggests new ways to treat Bnodes
  - By the algebra

# Counting

- What is a redundant answer?
  - BNodes in answers can be tricky
    - Several *notions* of redundancy
    - Distinguish between redundancy due to algebra and stated redundancy and inferred redundancy
  - Equality can be tricky
    - If two answers differ only by the value of one binding, and those values are inferred to be sameAs, how many answers?
    - No UNA in RDF-OWL
- We could count **answers** (instead of entities)
  - But then answers proliferate, often pointlessly

# Query Variable Position

- “Conjunctive ABox queries”
  - Standard in DL systems: KAON2, Racer, Pellet
  - No variables in property or class positions
- “Higher order” queries
  - `?x rdf:type ?C. ?C rdfs:subClassOf ?D.`
  - Careful restrictions make this feasible
- “Syntax reflective” queries
  - `?s ?p ?o`, where `?p` can bind to e.g., `intersectionOf`
  - `?x a [a Restriction; someValuesFrom ?C]`
    - Only bind to asserted axioms? (essentially SPARQL/RDF)
- Variants of latter two coming (see OWLED)



# More General OWL Queries

- What should `?x rdf:type ?C` return?
  - Graph: `:bob rdf:type :Person`
  - Some possible mappings for `?C`:
    - `:Person`
    - `owl:Thing`
    - `unionOf (:Person, not :Person)`
    - Bnode problem writ large!
- Some hope (but very sketchy)
  - Concept matching and unification
  - Traditionally defined for logics without disjunction

# For more thoughts

- See my presentation “Sparqling Queries”
  - <http://www.cs.man.ac.uk/~bparsia/2006/row-tutorial/>
- Also, at OWLED 2007 SPARQL/OWL?
  - Just enough to cover and encourage current implementations
  - Syntax for Distinguished and Non-distinguished
    - No Semi-distinguished except perhaps for explicit bnodes
  - Mixed Tbox/ABox queries
    - I.e., liberalize variable position
    - But also perhaps define a profile for strict Abox query